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# Eventual consistency vs. strong consistency: Making the right choice in microservices

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#### Abstract

This paper discusses one of the most crucial decisions for microservices, which is between eventual consistency and strong consistency. It examines how every model influences the system's efficiency, stability, and usability. Eventual consistency is also good for availability and scalability because it uses asynchronous updates, which is recommended for social sites like Facebook and e-commerce sites like Amazon. However, it brings issues like transient data instabilities and conflicts in managerial schemes. On the other hand, strong consistency guarantees real-time data updates, ideal for companies dealing with critical areas such as finance and healthcare since data must meet set regulatory standards. Still, it causes high latencies and additional expenditures on the infrastructure. In theoretical approaches, namely the CAP and PACELC theorems, this paper explains various aspects of the tradeoffs between consistency, availability, and performance. Analyzing the real-world use cases, including Saga and Outbox patterns, the most important approaches to achieve consistency with the needed freedom in system design can be outlined. In such a context, hybrid models appear as potential strategies that define consistency levels according to data and related operations' sensitivity. Anticipated future work, another trend in algorithmic adaptability and dynamic selfsynchronization, contemplates new strategies for the persistence of consistency at the distributed application level. Matching consistency choices with business needs allows organizations to improve system performance, costs, and credibility. This article calls for more comprehensive and long-growth learning, effective testing, and iterative architectural planning while dealing with the intricacies of consistency models in the dynamic setting of microservices architecture, providing artists and developers with a thorough guide in the process.

**Keywords:** Eventual Consistency; Strong Consistency; Microservices Architecture; Distributed Systems; CAP Theorem; Consistency Models; Domain-Driven Design (DDD); Event-Driven Architecture

#### 1. Introduction

Contemporary software systems often incorporate distributed environments for high availability, reliability, and dynamic resource management. In this regard, microservices architectures have emerged as optimal solutions for disassembling monolithic applications into manageable services that can be developed, deployed, and maintained individually. However, one major problem is the consistency and integrity of the data stored in these dispersed components. The general intent of this article is to help architects/developers and technology decision-makers navigate through consistency models to understand eventual consistency and strong consistency. The first area focuses on general and specific theories, the former a conceptual review of the topic while the latter offers an overview and details on their relative practical applicability in the context of readership and real-life cases. The article defines the relationship between consistency strategies, system stability, and performance by describing the benefits, drawbacks, and applications of the consistency strategies.

Consistency here means a situation whereby the data stored at different nodes are made the same in the long run, not the strict and immediate consistency defined earlier. Because it is an asynchronous solution, it dramatically improves

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system availability and capacity when network splits or high latency is common. On the other hand, this approach might lead to temporary inconsistency and, therefore, must contain reliable inconsistency resolution mechanisms for data accuracy. In contrast, SCA pursues strong consistency, which mandates that all data are updated simultaneously and either by a majority or via transactional delegation. While this model provides better predictability and more straightforward reasoning about the current system situations, it may hamper performance because of latency and lower reliability. Comparing these two paradigms, experts start to reveal how they affect design decisions in distributed microservices.

Microservices promote enhanced and flexible deployment, continuous integration, and self-sufficient growth. However, the same things become a weakness if each service takes different approaches to data updates or suffers from temporary disruptions. In user-oriented systems, including eCommerce, social media, and financial applications, unstable data states harm user trust and overall experience. For instance, an organization's failure to promptly update an order status or account balance would lead to confusion, duplicated transactions, or the complete loss of a customer's confidence. Sustaining data quality helps manage such risks since it provides a predictable operational context and envisioned effort at error control. Consequently, it enhances the manifold organization of several services, enabling some cohorts to introduce new accessories or expand previously existing ones without jeopardizing primary data quality.

This paper will provide a systematic approach to categorizing the types of consistency models and how they can be implemented in today's microservices architecture. It outlines both the theoretical framework and concrete examples of Sentiment Analysis, allowing the readers to go beyond abstract and basic concepts to see real-life applications. Such is the focus when comparing the CAP Theorem, illustrating how each solution promotes different business objectives, extending to the PACELC framework and newer trends in adaptation, such as AI-assisted adaptive consistency. In addition, the content also covers the psychological and user experience aspects of data syncing, as lack of synchronization may erode trust. The reader will be given lessons on reducing these effects via effective interface design, DDSD, and event orientations. The paper also becomes the manual that allows the teams of professionals to have the right knowledge they need to make accurate decisions that incorporate efficiency, expansiveness, and quality data. On this basis, the article creates the premises for further development of the problem.

## 2. Historical Evolution of Consistency Models

## 2.1. Early Distributed Systems and the Emergence of Consistency Challenges

In the early years of distributed computing, the main approaches concerned resource-sharing among nodes for enhanced scalability and dependability. However, as the use of these systems grew, issues regarding the consistency of their setup soon arose on large scales. Some nodes had to process requests parallel, and all had to have a consistent underlying data model. Due to the absence of elaborate replication methods, frequent problems such as stale reads, lost updates, and conflicting operations were identified. These challenges highlighted the need for consistency models to help manage updated commands across geographically dispersed nodes to enhance availability and performance. Hence, researchers and practitioners started to make definitions and structures with data consistency, which opened the first frameworks of theoretical conceptions that later prepared for generations of distributed system architectures and designs (Vogels, 2009).



Figure 1 Consistency Model in Distributed System

### 2.2. Key Milestones in Consistency Theory

CAP was one of the first and most significant consistency theories established by Eric Brewer in 2000. Essentially, the theorem states that at any given moment, a distributed system can only satisfy two of the three Cs: Consistency, Availability, and Partition tolerance. In practical terms, systems must commit to first order on properties when network partitions occur or to temporary unrecoverable failure. This led system designers to distinguish which of the two features – availability or strong consistency – was more important for a given application.

Soon afterward, scientists noticed that although highly remarkable, the CAP theorem did not account for the differences in latency in consistency decisions. Daniel Abadi (2012) introduced the PACELC theorem, which extends CAP by adding dimension; when there is no partition (P), systems must select between latency (L) and consistency (C). As a result, PACELC emphasizes that, other things equal, systems may need to maintain performance at the expense of data consistency. This added layer of awareness told architects that consistency trade-offs are not just seen when partitions are being implemented; they also apply in normal working conditions. It was beneficial to be more deliberate about these subtleties to improve the conversation about when and how often it is worth sacrificing other variables, like user-perceived latency, for immediate consistency.

#### 2.3. Rise of Multi-Primary Replication Protocols

When distributed applications were developed more, there was more pressure towards getting beyond single primary replication, where all updates come through a single leader node. Multi-primary replication protocols were developed to serve scenarios that involved multiple nodes accepting writes at the same time. In multi-primary systems, updates can be initiated at any node, which makes it highly available and has a throughput (Gray & Lamport, 2006). However, extending multiple primaries increases the challenge of achieving consistency because it is likely that one node will write data that another node also writes at the same time. It must resolve conflicts, which might be achieved by using version vectors or consensus algorithms, to ensure that all replicas go to the correct state. Multi-primary solutions enable increased parallelism, but they also require well-designed mechanisms to address the inevitable synchronization issues when implementing systems with this type of architecture.

#### 2.4. Adoption of Hybrid Consistency Models

Striving to achieve strong and, at the same time, eventual consistency predetermined the introduction of mixed strategies. Hybrid models permit system elements to use levels of consistency appropriate for the data and operations conducted on data sets. For example, a microservices-based e-commerce application may ensure that it behaves strongly for an inventory change but may be ok under eventual consistency for recommendations. This selective approach is advantageous when some data elements are more important than others (Gilbert & Lynch, 2002). Hybrid models always use techniques, such as conflict-free replicated data types (CRDT) or quorum writing, to ensure that low latency operations are managed with low consistency levels. This focused flexibility has made hybrid consistency viable for organizations that demand high performance and dependable data consistency across multiple subsystems.



Figure 2 Benefits of Hybrid Distribution Models

### 2.5. AI-Powered Adaptive Consistency

New trends have brought about the concepts of adaptive consistency that manage and regulate the consistency of AI and machine learning algorithms according to the current conditions of the real world. These systems can also help to

adjust the consistency demands during extremely high levels of traffic while increasing system responsiveness based on the observed workload, network delays, and user activities. During normal conditions, optimal consistency settings may be applied, as well as tighter ones when outside conditions are worse. The above adaptive mechanisms correspond with the essence derived from CAP and PACELC since they facilitate the barter of availability, latency, and consistency depending on the circumstances. But that is already a notable improvement that current, still in its infancy, AI-based techniques offer – an extra dimension of automation and intelligence. This innovation paves the way for other malleable architectures in the event of variations in workload.

## 3. How Eventual and Strong Consistency Fit into Microservices Architecture

Microservices architecture has become a fundamental shift in how large and complex software applications are designed in that they are composed of smaller, independently deployable microservices (Stonebraker & Çetintemel, 2005). Every service targets a single business capability, making it easier for teams to build, run, and extend specific tasks without affecting other system parts (Fox & Brewer, 1999). This approach has developed popularity among organizations looking to improve flexibility and modifiability in centralized settings (Birman & Joseph, 1987). However, since data is split over many services, maintaining consistency is critical (Lamport, 2019). Understanding how these two types of consistency fit with microservices will allow architects to pick out strategies that meet performance needs, reliability, and flexibility at the same time (Nyati, 2018).

### 3.1. Defining Microservices and Their Core Principles

Microservices are based on encapsulating certain business processes into individual services that interact utilizing simple contract interfaces. Each service also has data storage and management rules, which minimize the business-related overheads inevitable in monoliths. These are service autonomy, governance decentralization, and the essence of domain-driven design. With these principles, microservices help speed up the development and deployment cycle. However, such freedom triggers data synchronization issues, which means there ought to be well-defined plans about how synchronization will be handled. Much work is always being done to arrive at a consensus on where the service boundaries are, and data ownership has to be well-defined. This way, communication between different services can be kept to a minimum, and every service can employ the most consistent mechanisms that are adequate for its uses.



## 3.2. The Role of Loose Coupling and Service Autonomy

Figure 3 Loose Coupling Autonomy in Microservices

Inflexible coordination refers to the model of service provision where services are run with as little reliance on one another as possible so each can change at its own pace. This design promotes robust architectures because the downfall of one region is not transmitted system-wide. However, in data consistency, loose coupling exposes each service to

handle update settings independently. Services often need to use techniques that guarantee strong concurrency (synchronous mechanisms) or techniques that provide interest for eventual concurrency (asynchronous replication). In reality, there could also be highly consistent services on the mission-critical, which would require higher consistency to protect data. At the same time, there could also be lowly consistent services to enhance scalability. It creates liberation; each service can be constructed using dissimilar technology or paradigms. However, the architectural direction is still necessary to avoid cases of incompatible communication protocols or clashing data models.

#### 3.3. Why Consistency Becomes More Complex in Microservices

The consistency problem evolves when data is distributed across many domains, each of which may be handled by a separate service. The transactional models normally associated with monolithic architectures can become a problem in microservices. For example, payment, ledger, and notification services in a financial application all need to be updated frequently without the additional cost of distributed locks. Some of the services might be willing to tolerate temporal inconsistency to accept transient anomalies in exchange for quality of throughput, thus making them favor the approach of eventual consistency. Some may require force correction, leading them to good, consistent practices. Therefore, these models are used in conjunction with architects, together with the consistency strategies adapted according to the requirements of the domain. Hence, sound contract definitions around data, reliable messaging mechanisms, and clear monitoring are needed to identify synchronization problems. A synchronous/asynchronous hybrid mechanism can be employed to oversee the interactions between services that require high accuracy, and those reluctantly delayed slightly.

#### 3.4. Balancing Consistency and Scalability

Another elemental aim of microservices is to achieve horizontal scalability to accommodate many more transactions. High consistency is good as it affords correctness but inhibits scalability as it is a synchronous form of coordination. This approach may cause additional latency and higher contention in global environments, which is not the case in local environments. As for the type of data updates, eventual consistency uses asynchronous updates that usually improve liveness and availability. However, it creates temporary delta data that can interfere with user-oriented interactions, which are key to the performance. Most microservices-based systems use a mixed scenario where strong consistency is enforced on certain transactional operations, while eventual consistency is for other non-essential or less important operations. According to the ranked needs of every domain, consistent guarantees mean development teams can reduce resource overhead while addressing crucial concerns for businesses. Thus, microservices can be highly reliable and cover elasticity issues to remain effective with various loads. The decision and acquisition of the right consistency model are influenced by the organization's specified performance indicators, the data's importance and sensitivity, and the organization's capacity and willingness to handle more complexity.

## 4. Eventual Consistency Concepts

### 4.1. Core Principles of Eventual Consistency

Eventual consistency is a distributed systems model where all the data replicas converge to a particular consistency level even though they are not sent and updated immediately. These dissemination writes are asynchronous, contrary to writing where nodes always use strict write quorums for data, allowing for slight variations in replicas before they can agree on data accuracy. This approach considers availability and partition tolerance as the most significant factors, and based on the CAP theorem, extreme consistency cannot be achieved in large distributed structures (Vogels, 2009). When a system is in the eventual consistency, the opposite is ensured where, if time is allowed and no new updates are made, it will be seen that all replicas match the most current writes.

Synchronous updates ensure that the system runs continually despite the short-term unavailability of some nodes, thus reducing the likelihood of the complete system collapse. It can mean that the organization is inconsistent for brief moments, but it also means that such loose coupling is beneficial. Some authors pointed out that, in optimistic replication scenarios, conflict is only observed when replicas try to merge different versions of the same object (Saito & Shapiro, 2005). However, such a design usually requires resolution mechanisms for managing concurrent writes. Finally, the convergence principle guarantees that when updates are spread to replicas, a single consistent state of the distributed system will exist, encompassing all the valid writes.



Figure 4 Eventual Consistency in Distributed Systems

## 4.2. Examples in Microservices

Two typical microservices architecture approaches implement the concept of Eventual Consistency, namely Event Sourcing and Command Query Responsibility Segregation (CQRS). Event sourcing has its basis in the state of an entity in the form of one event, more preferably than a singular data record. As a bonus, this allows for recording the history of changes and making the change in a way that will be possible when updating other connected components asynchronously. Each time a new event happens, it is added to the event store, and the other services process these events at their own pace. Similar decoupling speaks of end consistency since each service handles updates independently without waiting for a global lock through wires.

Similarly, it isolates write actions from read actions of data, ensuring services have unique models for specific tasks. During a write, the command side has the update, while the query side reads from the denormalized or materialized view. This view is updated asynchronously, meaning there is a split-second delay between the most recent write and the query side. However, the approach allows the readers to go faster, and the contention is less since the primary dataset is not locked. With time, there is correspondence between the command and query models, representing the merging aspect at the heart of this architectural pattern.

### 4.3. Benefits of Eventual Consistency

The most significant advantage of the eventual consistency is its potential to ensure high availability regardless of the network conditions. Moreover, the system is elastic to regional blackouts or network splits by having each replica not necessarily be continuously online or in sync at the very instance when a write operation takes place. For example, Dynamo, developed at Amazon, is tuned to work well in networks that may not be dependable; it sacrifices consistency for availability (DeCandia et al., 2007). This design ethos enables microservices to stay responsive to business processes, thus enabling business continuity even when transient infrastructural breakdowns occur.

Another advantage is scalability because eventual consistency allows a distributed system to expand or decrease the number of replicas without the great cost of coordinating the new system. This characteristic aligns with the core of microservices architecture since services can be deployed independently of other existing services. Lax synchronization enables systems to fix loads that can accommodate many concurrent requests since strict synchronization reduces their effectiveness in handling them. Some of this elasticity is especially beneficial in unpredictable traffic scenarios that may see a sudden burst of traffic without consuming the optimum available resources. In conclusion, eventual consistency supports a robust and scalable microservices architecture, enabling the growth of the system while at the same time mitigating the costs of deprived strong consistency.

#### 4.4. Challenges of Eventual Consistency

The eventual consistency approach complicates the problem, and there are temporary inconsistencies in the data, which may be critical for some workloads. Several applications require time-sensitive, actual, and real-time data, such as those involving finance and medical fields. In such situations, the delay between the write process and the update completion for every node can cause issues or mistakes. Conflict resolution goes a notch higher, particularly when simultaneous write operations are across the different replicas. Whenever trickier merges of competing updates are needed, the necessary reconciliation strategies must be in place in such systems. However, these processes can get complex and have a high risk of error. Thus, partial staleness is an issue that developers have to examine carefully to determine if it meets their requirements.

The semantics of the user experience also require appropriate attention. Periodical delays in data occurrence can also harm users' confidence when there is a perceived inconsistency between application components. According to Abadi (2012), it is important to understand end-user operations when decisions are made consistently, resulting from design decisions for Database Management Systems, particularly where real-time data should be visible. Reducing confusion in microservices is usually achieved by providing signals that data may be stale or by ensuring that compensation strategies employed hide errors visible to the end user. Nevertheless, adding these features to application development can present higher complexity and development costs. Consequently, organizations have to find the right trade-off and implement high availability strategies while simultaneously considering that it might cause some level of user discontentment due to temporary freezes. However, detailed cost estimates have to be produced in the case of large-scale mission-critical scenarios.

Section	Key Points	Examples	Challenges	Benefits
Core Principles	<ul> <li>Asynchronous updates</li> <li>Replica convergence over time</li> </ul>	-	<ul> <li>Temporary data inconsistencies</li> <li>Conflict resolution complexity</li> </ul>	- High availability - Resilience
Examples in Microservices	- Event Sourcing: Stores state as a series of events	- Dynamo at Amazon - CQRS: Separate command and query models	- Lag between write and query visibility	- Fast reads - Reduced contention
Benefits	<ul> <li>High availability even during outages</li> <li>Scalability through relaxed synchronization requirements</li> </ul>	-	-	- Elastic workload distribution - Supports large-scale growth
Challenges	<ul> <li>Potential for brief inconsistency</li> <li>Complexity in conflict resolution</li> <li>User experience concerns</li> </ul>	<ul> <li>Financial systems (sensitive to errors)</li> <li>Healthcare applications</li> </ul>	- User trust may be impacted by data delays	- Allows operations to continue during regional outages
User Experience	- Delays might erode trust - Requires clear indicators or compensation techniques	-	- Complexity in mitigating user- facing issues	- Ensures responsiveness even under high traffic

Table 1 Key Aspects of Eventual Consistency in Microservices Architecture

### 5. Strong Consistency Concepts

#### 5.1. Core Principles of Strong Consistency

Concurrency consistency remains essential in distributed systems design, ensuring that any read operation maps to the latest executed write operation. This model maintains a setup whereby every node or replica has an identical view of the devices 'state without consideration of latency or network split. Strong consistency means the required components of a system are updated immediately, as all updates must be performed synchronously. A client observing the system state after a completed write operation will never have stale data; the system preserves a single view.

A significant component of high consistency is compliance with the ACID (Atomicity, Consistency, Isolation, and Durability) properties. Atomicity means that all transaction operations are completed or none is completed due to data integrity, and no half-baked data update can occur. When applied in this setting, consistency means that every change operation results in a state or the database that is 'correct'. Isolation helps transaction boundaries by isolating them from the side effects of other processing that may be occurring simultaneously. Last, Durability ensures that data after committing a transaction shall be permanent even if several failures exist. Altogether, four compose the basis of strong consistency and reduce new specific confusion over the state of data in distributed systems.

Concurrency also differentiates strong consistency from eventual consistency since it implies synchronous updates. Writing in the synchronous model can only commit the written value to all identified nodes before the system signals the commitment of the write operation (Katsarakis et al., 2020). This strategy minimizes cases where clients may read data that may have been updated a few years back. However, strongly constructed network communication and accurate coordination protocols affirm each update. While this strengthens a homogeneous data model, it brings temporal constraints that can be problematic in systems that cover different time zones or can temporarily or periodically lose connectivity.

#### 5.2. Examples in Microservices

In a microservices environment, the observed level of consistency is transactional, where a service gets updated in reference to correlated actions that are in a logical operation as a single unit. For instance, in an e-commerce application, a checkout operation involves a chain of services, including authorizing the payment, reducing stock, and confirming orders. In cases where the architecture necessitates this, strong consistency will be used to perform these services concurrently to avoid a situation where some operations are partially completed, leading to wrong inventory statistics or failed captures of payment information.

Another example is the ACID-compliant microservices patterns, where all of these patterns assume a common, strongly consistent datastore or one that coordinates transactions to adhere to unified records. This achieves the goal of the bounded context in that each service can be developed essentially in isolation. However, the system ensures consistent checks are performed across related services at transaction time. This approach is most essential when operations involve the exchange of monetary values or sensitive health information. Here, data anomalies are not acceptable, and correctness assurances are critical. High consistencies ensure that related components have the same view of business entities, thereby making concurrency, which may otherwise cause the complexity of reconciliation steps inaccessible.

#### 5.3. Benefits of Strong Consistency

The first apparent benefit of strong consistency is that it creates easy-to-implement application algorithms. Since developers assume that every read operation will reflect the last write, the application code will be simpler, and there is no need to worry about matching versions. The predictability offered by strong consistency supports clearer business rules: changes spread out right away, while transactions occur predictably. This is especially useful when governance or auditing is engaged or required since it delivers a simple and easily substantiated system state.

A strongly consistent system environment also provides high data integrity and user trust. Each user engaging in a service ensures that important procedures, including money exchange, are performed effectively and display up-to-date status. This certainty inspires confidence in applications that require the utmost accuracy and cannot afford to present incorrect data in real-life situations, such as online banking or electronic medical records. This strong consistency is equally applicable to many enterprise applications that require rigorous transaction control, which, in turn, increases the reliability of business processes. A key reason for the systems 'credibility is that users get the confirmation of their change requests as quickly as the results offer no perception of obsolescence by the users.



Figure 5 Strong Consistency in System Design

Strong consistency also reduces the chances of data inconsistency that may cause problems while implementing various services. When the data layer guarantees compliance with the ACID properties, concurrent writes are coordinated so that they do not lead to nonrecoverable states and multiple transactions. While contention issues may be present when performed in more general distributed environments, the nature of strong consistency everywhere provides a solidity blanket. This predictability benefits system administrators by making error recovery processes easy, where all the nodes have a consistent database view.

#### 5.4. Challenges of Strong Consistency

Strong consistency entails various significant challenges, which are described below. Latency and performance issues crop up because all nodes or different shards must agree before updates can be made. When the transaction is deployed in different regions or distant data centers, networks introduce delays in the commit time. As a result, some user-facing services could receive lower response rates. In some industries, they see a reduction in performance, which can affect user satisfaction even more. In highly competitive industries, architects must assess the value of immediate correctness against a possible loss in speed. Both availability and potential bottlenecks are critical factors for systems that require high levels of strong consistency. Since each write operation has to pass through several nodes, the system becomes sensitive to failures even within a single node or channel of communication. An outage of a node may cause writes to fail and cause availability problems if the cluster cannot quickly reroute transactions. When microservices are scaled, the overhead of used transactions, control, and coordination becomes a problem as it consumes more resources regarding network bandwidth and processing power. Organizations must provision infrastructure appropriately to ensure synchronous commits while not reducing the overall throughput.

The microservices architecture characterized by strong consistency can also be challenging to implement. The essential issues include, for example, coordinating distributed transactions, keeping the ACID properties, and avoiding deadlocks, which are challenging and need specific knowledge regarding interaction schemes. It could also need distributed locking mechanisms and complex concurrency control approaches, especially in cases of heavy traffic, which will increase the operational cost even further. This also applies to how architects design fallback procedures suitable if the network fails or nodes fail catastrophically. In some scenarios, a mixed model may be decided where only the most stringent applications enforce strict consistency, and in others, a more relaxed one based on the eventual one, though a less rigid one, will help to gain performance.

Accurate consistency remains a 'must-have 'where precision is crucial and has not been supplanted by weaker consistency variants. Synchronous updates and following ACID properties ensure uniform data visibility and strict compliance with needed high integrity in microservices. The trade-off is focused on the extra load hosted on system throughput and handling for errors that force teams to invest time weighing if the drawback of comprehensive, exact, and real-time reads outweighs the value possession to services 'pace. Strong consistency becomes most suitable when dealing with any situation requiring key business interaction or heavily reliant data reliability.



Figure 6 Strong vs. Eventual Consistency in System Design

## 6. Factors Influencing the Choice between Eventual and Strong Consistency

### 6.1. Nature of the Business Domain

The decision on which level of consistency will be achieved – eventual or strong – depends on the needs of the business domain in question. Data, accuracy, and compliance are requirements that often presume higher levels of consistency in financial systems. For instance, any transactional problems in a banking application would easily result in significant regulatory penalties if account balance integrity is involved. Like payment gateways, they need assurances that have been set to ensure that a particular transaction has been committed or backed out and reduce the possibility of double spending. Furthermore, industries with high regulatory requirements, including healthcare and insurance, require strong auditing, made possible through consistency and data traceability.

On the other hand, social media applications require frequent updates and embrace many users, enabling them to afford small temporary data inconsistencies. In these use cases, the users can wait a few seconds before seeing the likes, comments, or posts while the system gets to the final set. This tolerance is well explained by the fact that a delay in the pace of developing social contacts does not pose a threat to financial security and well-being. Therefore, the ones that care most about end-user interaction, often at the expense of data absolute accuracy, are most likely to gravitate towards this solution. When selecting consistency, the organization should mirror organizational needs with technology to ensure that the chosen model aligns with the organization's focal areas without unnecessary expenditures and complications (Kumar, 2019).

### 6.2. Performance Requirements

Concerning throughput, latency, and availability, the performance expectations act as determining agents of consistency. In cases where applications have to perform large numbers of writes and reads, even the introduction of synchronization in replication can cause performance degradation (Katari, 2020). Horizontal partitioning as a method to raise the throughput of the distributed database may be used in some cases, but always enforcing strict partition-level consistency necessary for achieving high availability may lead to increased latency. Therefore, high concurrency systems benefit from eventually consistent, highly available architectures that enable partial or delayed synchronization to release resources to run at optimum speed.

There are some scenarios in mission-critical applications where low latency has to be traded off against resilience. For instance, the real-time analytics stack uses hybrid consistency, where strong consistency is used on the most important data pipelines, with eventual consistency used for aggregated or historical data. These architectures use asynchronous replica updates to deal with burst loads and synchronous operations to protect the most critical transactions (Parker & Thomas, 2018). As a result, the gain on replication overhead has to be measured against the loss on data freshness. The best choice often depends on the relationship between the changes the business needs for operation and the state of stale or conflicting data.

#### 6.3. Tolerance for Eventual vs. Strong Guarantees in User Experience

The applicability of strong or eventual consistencies depends on user experience considerations. When data accuracy is critical to user confidence, achieving a strongly consistent option is necessary, especially during measures like an ecommerce checkout or airline ticket booking. A temporary condition of overlay may cause an understock of seats, leading to overbooking and hence negative experiences for the actual customers, besides producing wrong signals that a carrier is very popular and damaging its reputation. The error recovery strategies typically used in these domains include rolling back operations and user interface prompts; thus, there is a need to establish synchronous validity checks.

On the other hand, applications with mostly updates for non-critical data like content recommendations or friends lists can take advantage of eventual consistency with little complaint. People rarely need perfect coordination of the materials suggested by the app. Small inaccuracies, such as synchronizing suggestion lists a little while ago, are insignificant (Gilbert & Lynch, 2002). The designers in such scenarios employ background processes that synchronize data across the nodes. These processes can quickly regain most of the tension issues that would otherwise disrupt the processes. In other words, how much users can tolerate or even observe data inconsistency determines the feasibility of eventually deploying consistent state systems.

Factor	Eventual Consistency Considerations	Strong Consistency Considerations	
Nature of the Business Domain	- Suitable for domains prioritizing user engagement over absolute data accuracy (e.g., social media).	- Essential for domains requiring precise data integrity and compliance (e.g., financial systems, healthcare).	
Performance Requirements	- Benefits applications needing high throughput and low latency by allowing asynchronous updates.	- May introduce higher latency due to synchronous replication, potentially impacting performance in high- concurrency environments.	
User Experience Tolerance	- Acceptable for non-critical data updates where minor delays or inconsistencies do not significantly impact user trust (e.g., content recommendations).	- Necessary for critical data updates where inconsistencies can lead to user mistrust or operational issues (e.g., e- commerce checkouts).	
Technical Considerations: Database Choices	- Often aligns with NoSQL databases that support horizontal scaling and flexible replication (e.g., DynamoDB, Cassandra).	- Typically aligns with relational databases that enforce ACID properties and require strict transactional integrity (e.g., PostgreSQL).	
Sharding, Partitioning, and Replication	- Facilitates easier sharding and partitioning without complex locking mechanisms, enhancing scalability and fault tolerance.	- Requires sophisticated mechanisms like two-phase commits to maintain consistency across shards, potentially complicating scalability.	
Psychological and UX Perspectives	- UI strategies can mitigate perceived delays, such as progress indicators or notifications about data synchronization.	- Ensures immediate data accuracy, reinforcing user trust through consistent and reliable interfaces without noticeable delays.	

Table 2 Factors Influencing the Choice between Eventual and Strong Consistency

### 6.4. Technical Considerations: Database Choices

This is seen from the fact that the underpinning database technology does impact consistency strategies. In contrast, relational databases assure strong consistency since they support the commitment of resulting ACID–fulfilling transactions for cases that require the elementary notion of transactions (Marinho, 2020). However, scaling relational databases horizontally is difficult when every shard or replica needs to be synchronized at every moment. This increases developers' challenges, forcing them to develop features such as the two-phase commits that can be latency-inducing.

Many NoSQL systems intentionally use events to provide high write throughput and fault tolerance (Vogels, 2009). These systems can be partitioned or shared much more easily without having to work with distributed locks. There might also be certain conflicts, but well-developed conflict-solving strategies and proper schema planning will minimize

the risks and guarantee that data merge is performed correctly during its lifetime (Stonebraker & Cetintemel, 2005). In addition, administrators can use replication techniques that allow for a trade-off between the delay and data availability. For instance, being a distributed DBMS, some NoSQL systems allow the strictness per query or table to be tunable. The above flexibility lets each microservice conform to domain consistency standards while avoiding overhead in areas where ultimate rigidity is not required.

### 6.5. Psychological and UX Perspectives

Besides technical factors, user perceptions occupy a significant proportion of the consistency model choice. Brief periods of slow-forming content or flickering of the brief displayed on the screen may be excusable as long as it is restored quickly. However, consistency, which poses severe problems, can be the squaring of the parry or even call into question the accuracy of prior data presented. In the case of financial dashboards, customers are expected to confirm a transaction's processing reliability immediately. The presence of any sign of imbalance that is out of sync with other balances can cause skepticism about the reliability of the whole system.

To avoid associating UI with delays or inaccuracies, designers often incorporate UI strategies when they add reactions to a Web interface, for instance, adding progress bars or using disclaimers that inform users that a given piece of data may take time to update (Aguilera et al., 2007). Notifications can also clear up that the system is re-synchronizing or verifying changes. Thus, by informing the end-users about the presence of transient states before initiating the process of synchronizing, companies make the users more patient and confident during the entire procedure. This approach aligns with the principle of constructing transparent interaction to prevent forfeits of credibility, whereby user anxiety could be heightened in high-traffic events. To avoid letting the occurrence of eventual consistency turn into negative experiences for users, interface design is done in a way that optimally manages the expectations of the users. At the same time, effective error handling is incorporated.

## 7. Role of Domain-Driven Design (DDD) and Event-Driven Architecture

### 7.1. DDD Building Blocks and Bounded Contexts

Domain-Driven Design (DDD) is a strategic approach to creating complex software systems whose essential components are based on specific domains in real life. One of the key principles in DDD is the bounded context, which reflects that a model should contain only elements pertinent to a given functional area. Applying ubiquitous language to teams means consistent terms will help communication between domain specialists and developers. The benefits include the ability to identify changes separately and simplicity because a specific BC comprises only a subset of the business logic. Besides, aggregates are consistency boundaries because if two users want to write, they can do it in realms of defined entities. In the same manner, Evans (2003) also explained that these constructs help to clarify the communication between various organizational units because each context can be modified separately within a system without subjecting the system to the risk of having all the units supplying inconsistent information. This way, domain models are divided into separate modules, making teams more gifted in the kinds of design they are willing to inculcate; this way, core business objectives are met, Johns 2005). This sub-dividing, coupled with a well-crafted context demarcation, forms the staple of most of the accomplishments documented in DDD practices.



Figure 7 The building blocks of domain-driven design (DDD)

#### 1.1 How Event-Driven Architectures Leverage Eventual Consistency

By being event-driven, these architectures are strongly aligned with many of DDD's tenets in that they support decoupling and allow services to operate independently. As such, services transmit and receive events; therefore, this shift facilitates a more scalable and reactive approach to system design. Whenever an event causes a subsequent update, such updates do not reflect immediately on all the components, creating temporary inconsistencies. There is always a way of closing such gaps because, in the end, replication and reconciliation must cease, and the information must stabilize (Fowler 2003). This approach allows services to scale out and deal with increased load independently, minimizing bottlenecks. In addition, event-driven architectures enable non-synchronous processing, which means that tasks take different amounts of time to be done, and this can will not in any way cause a poor user experience. Decoupled services also bring the flexibility of failing independently of others. If a particular service fails, the remaining services will continue until the faulty one is back online. Therefore, an event-driven model helps teams propel high volumes of information while following consistent interest with workability.

#### 7.2. Aligning DDD Principles with Consistency Choices

In order to utilize DDD to the maximum potential in an event-driven architecture, the consistency preferences should be matched with the business needs being served by each context. There may be situations when some bounded contexts require strong consistency. At the same time, in other cases, it is possible to use eventual consistency if it will improve the system's performance and extensibility (Taylor et al., 2010). For example, a context managing financial flow can have ACID transactions to ensure data integrity. At the same time, other services, such as notifications, could be updated in the background. According to Vernon (2016), this fine-grained aim and consequence design lets architects determine conditions under which and how synchronization occurs while avoiding unnecessary coupling of lowpriority features. This also helps minimize the use of various resources and simultaneously protects key system functionalities to be reliable. While aligning the domain model with an appropriate level of consistency, optimizing the flow of events, defining adequate margins, and creating smooth communication between services is possible. In the end, this alignment makes it possible for clients to realize the intended business goals of each context whilst maintaining the architectural independence that underlines DDD.

### 8. Guidelines for Evaluating and Implementing Consistency Models

When choosing consistency models in microservice applications, one must balance the need for consistency with business objectives, technology limitations, and core processes. From here, the notion of a consistency level can be mapped by architects to the demand of such service and the data class. However, the distributed computing paradigm's real challenges can be reduced by exposure and testing procedures. The following guidelines present a step-by-step approach for assessing and deploying consistency models. They are simplified concerning their original presentation, based on rigorous research and practical experience available in the academic and industry literature.

### 8.1. Conducting a Requirements Analysis

The first step in obtaining a detailed requirement analysis is distinguishing between operational and non-operational data. Privileged data, like financial operations and metrics important for operational and tactical management decision-making, typically necessitate high levels of consistency to ensure data integrity and user reliability. Small discrepancies in these areas threaten regulation or reputation, so more consistency assurances are needed. On the other hand, auxiliary data, which may include unrelated user preferences, can sustain eventual consistency without much impact on application usability or user appreciation.

Data has to be classified, after which correct latencies are to be decided, and service level agreements (SLA) need to be established. Real-time systems that can be found in high-performance requirements like real-time bidding can require responses in sub-second time, which creates a need for high read rates while at the same time generating high update rates (Gorenflo, 2020). On the other hand, relaxed communication styles that allow batch processing can afford higher latencies at the cost of simpler synchronization. Through formal SLAs, it becomes possible to connect consistency objectives to performance expectations. The above process is supported by earlier work that assumes clarity in distributed system design.

Another component that influences consistency requirements is regulatory compliance. For example, financial and healthcare sector data accuracy and audibility requirements require stronger consistency models, thereby excluding weaker models. By evaluating these constraints from the onset, the teams can define the right blend of consistency approaches. In various organizations, there are combined approaches in which tight control is exerted when it comes

to core missions brought about by consistency. At the same time, the less stringent models are used for additional services (Nyati, 2018).

#### 8.2. Architectural Considerations

Architectural decisions are next to address network topology and replica placement concerns once requirements become more definite. Although replicating data increases the availability and reliability of the data, it comes with the cost of increased latency and complexity if the system has to maintain a strong consistency over data. Proper choice of replication techniques like the leader-follower approach ensures the data remains consistent even under failure cases. Some replication models incorporated into the design of Dynamo, for instance, make use of eventual consistency to handle read and write intensity while simultaneously containing the degree of unavailability (DeCandia et al., 2007).

Both monitoring and observability come as essential categories of consistency management in transit. Having distributed tracers, centralized logs, and real-time panels makes it easy to identify disparities and preeminent latency. To give insight into the message queues, replication lag, and commit rates may point out situations where strong consistency decreases performance below some important threshold. On the other hand, it shows whether the system reconfiguration due to these anomalies is caused by eventual consistency. Strong monitoring aligns with industry standards when organizations focus on proactive fault identification (Vogels, 2009).

System architects must select the consistency model procedurally by matching it with the database options. Relational databases usually have transactional assurances but may compromise performance in fully distributed systems. Usually, much of the complexity and additional logic to deal with conflicts logically partitioned and tolerant is also present in NoSQL databases. In any environment, whether inside the company's facility with an on-premises infrastructure or in a public or private cloud with a platform, the physical placement of the servers, clusters, and network links determines the practicality of implementing some of these consistency levels.

Identifying those heavy loads may be useful for comparing the chosen consistent model's performance when demand increases. In certain circumstances, horizontal scalability in services may also lead to increased replicational lag, which forces reconsideration of existing structures. That type of information can be used to orchestrate change in architecture to sustain dependability and avoid earmuffs of degradation.

#### 8.3. Testing and Validation Approaches

Strict test and validation procedures to ascertain whether the selected consistency model is adequate for the organization. Chaos engineering that inverts the system to create a catastrophe can help identify potential issues with replication approaches, unveil a slow means for replication, or demonstrate data contention (Gray & Lamport, 2006). Regardless of the approach used, the effects of having different subsets of nodes or individual nodes fail or intentionally be separated can be used to quantify the time needed for the data to converge to an eventual state or the reliability funneling a strict consistency model affords.

Failover simulations constitute another testing category because they show the architecture's performance during hardware outages. These simulations check whether redundancy mechanisms work properly and whether replication techniques are invisible to the end-users. Observed user's time and logs of applications and data correctness during failovers provide an understanding of sacrificing high consistency for availability.

Staging environments also significantly contribute to the validation of consistency models. Such traffic resembles a production system's working conditions and enables testers to see how the system works without affecting actual customers. Staging environments are used with canary, release, blue-green deployment, or any other rollout to ensure that these new consistency configurations do not cause more breaks. These validations support a change-driven process based on performance data, user feedback, and other operational data to fine-tune consistency commitments.

Setting criteria for analyzing and applying consistency models in microservices architectures requires carefully and thoroughly examining data needs, architectural actions, and testing approaches. Understanding the difference between mission-critical and auxiliary data is instrumental in deciding when eventual or strong consistency is suitable. The references to the specific setting of the network topologies, the nature of replication, and the monitoring instruments contribute to the sustainability of the chosen model at scale. Testing is done through chaos engineering, failover testing, and staging environments that help build confidence in the given configuration. In the end, they foster a strong and forgiving organization that can respond to changes in business requirements. This paper will further demonstrate how incorporating principles of domain-driven design and adopting the 'growth mindset' of the culture of permanent improvement can help preserve data consistency in the context of a constantly emerging and expanding, large-scale

system of microservices. Therefore, when an organization follows these key guidelines, it is possible to achieve performance, availability, and data integrity that ensures user growth and trust.

### 9. Patterns and Anti-Patterns

Architects often employ predefined paradigms concerning how data consistency should be addressed in microservices and tend to avoid quirks that result in unsuitable circumstances. Some architectures include Outbox Pattern and Saga Pattern to allow organizations to live with eventual consistency where amplitude is part of their architecture, whereas strong consistency tends to use the Two-Phase Commit (2PC) or distributed transactions. As important is learning what not to do, which includes overusing the two-phase commit protocol and attempts at a monolithic data model, which are counterproductive to microservices.



Figure 8 Two Phase Commit Protocol (2PC) Process

### 9.1. Common Architectural Patterns Supporting Eventual Consistency

### 9.1.1 Saga Pattern

The Saga Pattern is often used at the application level to manage long-lived transactions involving multiple services to be managed consistently. While a saga transforms multiple resources into a single, long transaction that cannot be committed until all operations finish, a saga reduces each workflow step into a sequence performed by either a choreographer or an orchestrator. Services envisioned as participants in a dance interact through events; each receives these events and follows the proper compensation or forwarding steps. This design lessens tight coupling since every service is independent and invokes the subsequent step via an occasion emission. While sagas encourage availability and responsiveness, they can complicate the implementation process, particularly when designing compensating transactions for semi-failures (Gray, 1981).

### 9.1.2 Outbox Pattern

Consistent models of the outbox pattern deal with the problem of updating several services or data stores. The outbox table and outbox log can be used so that the event and the primary data update fall in the same transactional scope. The outbox in a post-commit update model stores information describing the specific event to be published for each changed data item. After a specific commit, another process reads from this outbox and publishes these outbox events to a messaging system, guaranteeing strong consistency between data changes and event outbox publication (Stonebraker & Rowe, 1986). Event emission is separated from the primary transactional path in this pattern, contributing to message loss prevention, idempotency, and scalability. The trade-off comes in the form of extra operational complexity for having two outbox processes to manage and avoid creating unduly large delivery latency.

Eventual consistency patterns can also shift the burden of timing for data convergence back into the hands of developers. Under certain circumstances, systems may be characterized by brief violations that a user can recognize. This can be solved in the Saga Pattern by compensating actions and in the Outbox Pattern by replaying events, which can get data into line after short periods. Lastly, the mentioned patterns demonstrate how microservices benefit from asynchronous communication and carefully designed workflows to remain highly available and with minimal locking overhead.

### 9.2. Common Architectural Patterns Supporting Strong Consistency

### 9.2.1 Two-Phase Commit (2PC)

The Two-Phase Commit base protocol must be used to ensure two-plus consistency between different services. In 2PC, a coordinator service coordinates and completes a transaction with every participant that either gets fully committed or rolled back. This option avoids partial updates and keeps data consistency and integrity at an optimum level (Bernstein & Goodman, 1983). While it provides a certain degree of transactional safety needed for crucial operations, 2PC can result in a high level of latency and may become a bottleneck. The coordinator also has to get approval from every participant, which imposes much load on the network, thus putting the whole system at risk in case of failure due to the coordinator's breakdown.

#### 9.2.2 Distributed Transactions

In addition to 2PC, distributed transactions are a high-consistency solution. This model invokes services under a global transactional boundary monitored by a distributed transaction coordinator. Using the ACID properties across the cluster, the manager guarantees that all the nodes see the same data,tate (Abadi, 2009). This can happen when an exact degree of correctness is needed, for example, in financial or healthcare systems. However, the distributed transactions could lead to the need to use special middleware, face complicated debugging, and lower throughput with the growth of services.

Strong consistency patterns remain essential when data consistency cannot be violated. In return, developers must consider higher operational complexity, performance expenses, and possible centralized issues against assured and verified data for microservices.

#### 9.3. Anti-Patterns That Lead to Poor Performance or Data Integrity Issues

#### 9.3.1 Overusing Two-Phase Commits

Although 2PC provides a strong consistency, relying heavily on this mechanism negatively impacts a microservices environment. The fact that each participant awaits the coordinator's green light exacerbates system-wide latency and diminishes its tolerance to node crashes. Some disadvantages of using 2PC include Resource locking, which makes scalability an issue. There are often cases where eventual consistency and compensating transactions are enough and even faster. If they over-apply 2PC, the teams create a highly-coupled system. When one part fails, it pulls down others and defeats the benefits of microservices.

#### 9.3.2 Monolithic Data Models in a Microservices Context

There is another significant anti-pattern when microservices use the same schema but are monolithic. Focus on centralized data structures contradicts the conception of complete SOAs' services independence: they stay interconnected. This also goes against independent deployment and presents problems of scaling specific modules without affecting the rest of the system. For large, unified data models, concurrency conflicts are even more difficult to resolve when services interact using the models. This approach goes against the DDD best practices, at least as far as the version proposed by Evans, because each Bounded Context should contain its data model, which should correspond to a particular functional area.

These designs also worsen the problems of versioning and could provide ways of distributing errors across many unrelated services. If one service changes the schema, all other services may be compromised or deliver data incorrectly. The decentralized data approach, on the other hand, is an advocate for isolation, maintainability, and scalar. Eliminating this anti-pattern means undermining unclear data ownership and ensuring each service takes its portion of the total domain (Esas, 2020). In this context, it is possible to introduce the patterns that will help organizations achieve the benefit of having microservices and that are more aligned with the philosophy of microservices. Choosing an eventual or, conversely, a strong consistency approach is possible depending on business specifications, the level of risk that can be absorbed, and performance limitations. However, there are anti-patterns, such as over-relying on 2PCs or having centralized data, which should commonly be looked out for and fixed to ensure health. Implementing patterns and anti-patterns in distributed environments requires evaluation, regular enforcement, and a readiness to adapt to the progressive aggravation of challenges as organizational environments grow progressively diverse and intricate.

## 10. Cost Implications

#### 10.1. Impact of Consistency Models on Infrastructure Costs

Whenever different consistency models are used in a distributed system, there is a tendency for the infrastructural costs to increase because of the need to consume more resources, especially where high levels of consistency are needed. High consistency often calls for synchronous operation between nodes, leading to many replication loads. For example, each write operation may contain several confirmation messages, resulting in higher CPU utilization rates and memory allocation. As such, increased costs relating to the procurement and management of server resources become an issue of concern in organizations. These cost increments become compounded as systems grow, affecting financial planning and reduction schemes.

Overhead costs, in turn, fall mostly under replication and storage costs. Extensive consistency requirements dictate that data be kept in check by ensuring that all copies of the data replicate the other, increasing storage requirements where data is copied often. As described by Stonebraker and Cetintemel (2005), when the degree of replication is raised while fault tolerance is augmented, operational intricacy increases as well. Likewise, the network bandwidth usage is also high since the nodes operate in a parallel fashion and send updates to other nodes to ensure consistency. These factors can be particularly critical where big data or geographically dispersed structures are within the system. This aligns with network bandwidth and latency costs that exert considerable pressure on total expenditure levels. It is acknowledged that the need for the regular synchronization round can occupy a significant amount of bandwidth, meaning that the utilization of services provided based on cloud computing alternatives will imply the necessity to consider the level of the ongoing costs as well. However, applications achieving strong consistency for latency-sensitive applications may need some extra technologies in the areas of networking to minimize communication time (Lamport, 2019). Although such measures can spur performance, they lead to higher monthly service fees, as seen above. Therefore, it becomes necessary for organizations to weigh the importance of having data instantly synchronized against available network budgets.

#### 10.2. Budgetary Considerations for Scalability and Performance

Scalability is an important cost driver as services requiring fast scaling may increase infrastructural costs associated with reliable consistency models. As the distributed system increases in size, achieving consensus globally can be costly, leading to organizations having to use the best hardware (Gray, 1976). The second approach, Eventual consistency, is relatively more cost-effective because it permits partial and eventual updates that do not require continuous synchronization. However, eventual consistency leads to conflicts, and this, in turn, may require more resources for the correction of those conflicts.



Figure 9 Budget scalability

The tradeoffs between lower cost and higher consistency are technical and business. Any systems dealing with important business transactions like electronic fund transfers normally demand much overhead to ensure the data is immediately consistent (Gill, 2018). In these cases, the potential for data imprecision can be expensive when it reaches the monetary dimension, hence the need for elaborate consensus protocols. Applications that do not require transactions, such as social networking services, could prefer delay-at-least-consistency over state-machine consistency to cut costs but meet acceptable client experience. However, effective management of these costs requires careful optimization of cloud deployments. Some organizations apply hybrid solutions, maintaining strong consistency in

sections of the architecture that process relevant financial data but permitting eventual consistency in the features that accept a certain amount of delay for their functions (Vogels, 2009). Companies can self-scale cheaply through the division of service based on consistency requirements. This practice also allows teams to choose a certain server type or instance size and plan the workload of a particular service and its availability.

Another common feature is the on-demand elasticity that enables the provision of services in proportion to their usage. It may minimize the risk of over-provisioning since it gives freedom to adjust the solutions provided to match the expressed demand but also results in cost variations. Consistency mechanisms that rely on recurrent, synchronous updates cost operational expenditures, especially when traffic densities arise. Researchers recommend decision makers pay close attention to key performance indicators related to application performance in the post-IPC environment and make necessary tradeoffs between the need to ensure a consistent end-to-end data view across the enterprise on the one hand and managing costs that could negatively impact the bottom-line on the other. Furthermore, the measures, including caching and read replica, can offset some costs by limiting write operations, whereby the bandwidth charges will be drastically governed. Deciding on an optimal consistency model involves managing infrastructure costs while maintaining efficiency and availability to achieve the highest levels of dependability for organizational applications at a reasonable cost.

## 11. Recommendations on Aligning Consistency Choices with Business Goals

### 11.1. Balancing Risk, Cost, and Performance

Risk, cost, and performance are the main trade-offs organizations must make when choosing the consistency to apply in their organization. While achieving higher degrees of consistency may be beneficial, this can result in higher infrastructure costs since more entities like distributed transaction managers or synchronous replication are typically needed to maintain consistent data copies across nodes (Morgan, 2015). However, occasional inconsistencies resulting from eventual consistency can reduce replication costs and bring in [(2)] better horizontal scalability. When risks are assessed, it becomes possible to understand which data operations require synchronous assurances and which processes can take delayed updates without much loss. Approaching it in such a way guarantees resource utilization efficiency, addressing business functions that benefit from stronger consistency. Moreover, the performance of implementation decisions needs to be integrated into the macro budget prerequisite, as overemphasizing highconsistency architectures could deter innovation outcomes that result when resources are allocated to other strategic domains. On the other hand, failure to observe relevant consistency assurance measures may lead to a loss of reputation or monetary losses for the business. Consequently, tolerance to risk, cost, and performance will define the overall organizational capacity to satisfy its operational workload demands.



11.2. Tailoring Consistency Requirements to Business-Critical Operations

Figure 10 Types of NoSQL Databases

Adopting differential consistency standards applicable to critical business operations is crucial for effectiveness and efficiency gains (Jones, 2017). Operations such as financial transactions require higher degrees of consistency lest wrong debits or credits occur, while real-time analytics can occasionally allow lower consistency if it means faster

processing. This partitioned approach is coupled with the domain-driven design, which breaks down a system into bound contexts requiring different data accuracy levels. Thus, architects can separate the services that need real transactional guarantees and implement consistent mechanisms without overloading the entire system. Moreover, identifying appropriate databases for use with relational or NoSQL technologies means that the chosen models are compatible with the technology's inherent characteristics (Lee, 2019). Also important to mention in this respect is the definition of the service-level agreements (SLAs), which define when latency is acceptable and when systems downtime is okay. These SLAs help to identify to what extent consistency measures are still successful as special precautions against changes in load. Finally, when determining combination consistency and considering the criticality of microservices, organizations occupy an optimal balance of throughput-to-integrity and guard against consuming too much operational effort.

### 11.3. Continuous Review and Iteration

Consistency choices should be made with requisite flexibility because review and iteration are critical to achieving and sustaining alignment on business goals. Continuous audits ensure that organizations affirm whether the selected models remain effective in sustaining the other measurement parameters, including system throughput, response time, and fault tolerance (Green, 2016). This is because automated monitoring can easily detect any source of congestion or high latency, which can be corrected. In addition, stakeholder feedback indicates the changes in customer expectations and the new requirements set by the regulators. As these factors evolve, there may be a need to switch between microservices, some of which should be considered as altering the consistency level from a strong consistency perspective to ensure that data integrity and user satisfaction are maintained. Documenting each modification helps prevent confusion and makes it easier next time someone wants to audit the process for improvement. This adaptive process plays a vital role in eliminating extreme views of organizations and performance objectives by balancing them while keeping the consistent models selected in a strategic line of sight. The iterative approach ensures that scalability and cost advantages are maintained and consumers are protected.

## 12. Case Studies or Hypothetical Scenarios

### 12.1. Financial Services Application with Strict Transactional Needs

People who use financial organizations need ironclad guarantees to protect accounts and transaction data. For example, for a micro service architecture for a global bank with various services across multiple regional centers, there needs to be a very high level of consistency in real-time deposits and withdrawals so that there is a true reflection of the financial status of a client. An architecture involving synchronous communication and the ACID properties is implemented to avoid any partial updates affecting data integrity. A two-phase commit protocol is often used to enable transactions across many services, but it raises latency levels in the process. This design helps ensure that each entry is easily traceable and meets local rules and regulations while lowering the risk of expensive audit failures. This approach is fully justified, given that even a single discrepancy might entail negative consequences in the industry's regulation. Possible threats such as double spending are also mitigated by the constant validation of transactions, resulting in the application of Gray and Lamport (2006), according to which 'transaction commit and succeeding outcome require consensus'; similarly, the reliability highlighted in this model becomes accrediting when evaluation is made of its significance in financial conditions.

### 12.2. E-Commerce Platform Handling High Traffic with Partial Eventual Consistency

In the large scale of e-commerce, microservices handle conditions such as catalogs, orders, and user profiles when demand is unstable. A combination of frugal and eventually consistent access methods is realized to ensure efficiency during the relatively more intense shopping time. Operational parameters, such as the update interval of inventory details, incorporate asynchronous replication, making most query wait time minimal. While the firm's critical data services processes, like order confirmation and information consistency checks, are tight and rigorous, other, less sensitive services incorporate loose and friendly consistency to enhance responsiveness. This combination of rigorous and progressive consistency assists in relative mass transaction issues because the system is graceful instead of wholly failing under great pressure. Application caches are added to buffer backend loads, meaning that caching tiers are added to handle spikes in consumer traffic while sparing transactional capabilities. This pattern supports Bernstein et al. (1983) guideline to maintain transaction-middleware equilibrium; thus, supporting flowing and visible transactions if being disrupted temporarily does not severely harm user confidence.



Figure 11 E-Commerce Website Scale-Up Strategies

### 12.3. Social Media App Prioritizing Availability over Immediate Consistency

A social networking website that supports millions of daily active users inevitably focuses more on its activity and timely feature releases rather than unvarying stability. Within this framework, the rate of data updates, for example, real-time messaging, might be more tightly regulated, whereas features like social feeds, likes, or comments might be easily delayed. Eventual consistency models allow the back end to replicate data changes among conjugate nodes without requesting simultaneous acknowledgment. The architecture incorporates replicated services across many regions, so if a particular node experiences latency, a user request is immediately redirected to an alternate site. Elements like push notifications and content timelines are fast and do not sap the positive user experience, even if a certain degree of synchronicity suffers from milli-second inexactness. The approach aligns with Taylor et al. (2010) categorization of how intercoupled service interactions promote highly scalable designs. The users are not immediately aware of small content delays and can be content with staying connected while the platform remains operational, avoiding huge numbers of disconnections during network splits.

### 12.4. Lessons Learned and Best Practices

Different needs arise for strong or partial eventual consistency patterns to be implemented in all these usages. Outcomes from the financial services leakage case focus on guarantees in highly sensitive sectors where data errors attract stern consequences. On the other hand, the context of e-commerce explanations shows that an approach of strong consistency in strictly necessary cases while using eventual consistency in other fictional cases makes sense. The social media scenario exemplifies how user-space features can benefit from weak consistency as long as availability is a priority at all costs to the business. An examination of key decision trade-offs shows that there are fundamental compromises related to consistency models, which include needing to be in harmony with operational goals and objectives, as well as expectations of the users. As mentioned by Krishnamurthy and Rexford in 2001, these distributed applications should periodically collect performance information to support the dynamic adaption of data copying and the flow of application services. Another advantage of incremental deployment is that organizations can test each consistency mechanism's impact on the latency and throughput level before deploying the strategy across the entire organization. Enhanced logging and the use of analytic measurement in real-time add to the readiness for immediate identification of abnormalities to encourage more preventive action rather than reactive. These are also supported by industry experiences where slight changes in consistency models affect SLAs and the architecture of error-handling mechanisms. Ideally, solutions can be retained 'close to the chest' AND implemented via a combined effort of development, operations, and business teams. Consistent decisions based on informed data, on top of which monitoring is also implemented, allow for achieving reliability and scalability objectives and responding to new requirements.

## 13. Emerging Trends in Distributed Systems and Consistency Models

### 13.1. AI and ML-Powered Adaptive Consistency

Distributed systems are, therefore, able to experience adaptability and consistency by introducing AI and ML. Considering the frequency of users 'requests, network delays, and data access rates, the consistency values can be set dynamically by utilizing the ML algorithms, depending on the current priorities, such as performance or increased reliability. This is unlike static topologies because it enables the system to adapt to the ever-changing load, failure of some hardware, or increase in traffic. The author Vogels (2009) also states that new techniques can be designed that

allow auto-tuning of consistency based on feedback mechanisms that can indicate when consistency levels can be relaxed or where they must be rigid. This requires various adaptive techniques that apply real data to foresee uneven performance before it affects users.

#### 13.2. Dynamic Scaling of Consistency Levels

Equally important, the general idea of dynamic scaling of consistency levels has also been receiving growing interest. Therefore, prior solutions for distributed systems are usually designed based on a single consistency model and cannot easily accommodate dynamic workloads. Recent studies indicate that it is possible to gain great performance advantages with algorithms that offer stricter guarantees of consistency only for transactions requiring it and more relaxed pledges for transactions that need not be quite rigorous (Tanenbaum & Van Steen, 2017). This approach presents a need to complement the decomposition of application components with constraints under which such parts can manipulate data. For example, financial transactions may be performed at a higher isolation level than read operations, which in turn may use the techniques of eventual consistency in order to be as fast as possible. These are also supported by nonblocking commit protocols that enable partial transaction progress even when a single node is experiencing communication lag, as elucidated by Skeen (1982). In other words, dynamic scaling strategies enable organizations to adapt effectively to balance between accuracy and processing capability.

#### 13.3. Impact on Future Architecture Designs

Regarding the coming architectural visions of other experts, they also expect AI-driven, dynamically scaled consistency strategies to define the future state of distributed systems. At the same time, an image of improved observability schemes is developed that record metrics at various levels, including the network level and interactions between applications (Fox & Brewer, 1999). With such data concomitant, architects can better understand how consistency settings affect user results. Improved decision support tools are anticipated to provide the capability for automating recommendations and identifying any sign of deviation from the optimum performance level in the quickest time possible. However, as more organizations adopt edge computing products, the challenge of correctly synchronizing the data received from dispersed geographical locations increases. It is possible that using adaptive consistency mechanisms enhances the application of AI by learning local access patterns and easing constraints whenever feasible. In the long run, the distribution system is a spectrum of customizable, adaptive, and automatic consistency based on real-time analysis and predictive models. While using availability versus consistency tradeoffs, developers and architects will leverage integrated and data-driven feedback loops to improve the resilience and performance of service at scale.

In addition, when container orchestration platforms develop, microservice applications are available to be deployed in their usage consistency of real-time (Gogouvitis et al., 2020). This shift highlights the need for intelligent load balancing in which monitoring engines powered by Machine Learning algorithms seek to balance the highest level of data accuracy against reasonable response times. In the coming years, enhancements will be made to the models by academic and industry collaboration, which will bring down overheads by percent, without any loss in precision. The next generation of distributed systems will continue to be massively agile and extremely robust in the face of varying demand levels, with next practices borrowed from concurrency control, replication, and advanced analytics.

### 14. How Advancements in Tools and Frameworks Might Impact Decision-Making

### 14.1. Global Consistency Solutions and Their Trade-Offs

With global databases' developing availability, architects' ways to maintain consistency in such microservices have become a critical concern. Services like Google Spanner, CockroachDB, DynamoDB, YugabyteDB, TiDB, FaunaDB, and Amazon Aurora Global Database have different ways of maintaining consistency across nodes in different regions. For example, Google Spanner applies a globally synchronized clock to maintain external consistency, reducing complicated data copy divergence and simplifying application logic for programmers. Similarly, CockroachDB is a distributed, highly consistent SQL that uses a raft approach to deal with replicas. However, these make the systems have higher latencies because of the overheads accompanying synchronous replication.

Typically a NoSQL database, DynamoDB provides the option to enable configurable levels of consistency. Using this strategy, developers can have total control over the strength or eventuality of each request, aiming at tasks that need variable consistency levels, both low-latency reads and ordering (Carbone et al., 2020). On the other hand, YugabyteDB incorporates a distributed SQL layer of high-speed replication into the plan to merge relational operations with well-being construction. These systems help minimize the overhead of operations. However, they do require an optimum

level of effort in what schema design, network setting, and resources would keep the system at the right level of consistency.



Figure 12 YugabyteDB Architecture

TiDB and FaunaDB focus on horizontal scalability and transaction consistency while implementing replication strategies in different ways. TiDB scales computation and storage independently and aims to be a globally distributed serverless database that offers strong consistency like FaunaDB. Nevertheless, this architecture can imply the emergence of communication complications at nodes in any distributed system. The use of distributed architecture in this regard can pose problems of clock synchronization, as highlighted by Lamport (2019), which creates doubts over the ordering of transactions when integrated into a given system, a factor that, in our view, suggests that their function must thus be given a rigorous performance review to determine suitability.

A few organizations use a mix of these methods for critical and non-critical data, which may indicate an added level of complexity. Implementing the deployment of such systems requires extra attention to overcome transactional integrity problems during both normal and exigent circumstances.

### 14.2. The Growing Ecosystem of Managed Services

In addition to globally distributed databases, the notion of managed services has become more prominent, which hides much of the provisioning, maintenance, and scaling of consistency-focused infrastructure. Others include Amazon RDS, Google Cloud SQL, and Azure SQL Database; they relieve organizations from many simple tasks, including but not limited to patching, replication setup, and failure recovery. This managed approach proves more advantageous for teams whose major concern is not with the database but with business logic. However, the migration decision that can lead to adopting managed services raises questions on vendor lock-in, service-level agreements, and integration with on-premise or other clouds (Balobaid, 2020).

This preference for managed services can be attributed to one key factor as enterprises gravitate towards their cloudnative counterparts: the ability to remain resilient and flexible. By this, a managed solution can minimize the complexity of compliance by backing up data and encrypting data at rest by monitoring data in real time. However, some qualifications are required to determine whether the choices offered by the service meet the needs of a particular application. In cases where live data consistency is required, some of the managed offerings might set up the read replicas in such a way as to include propagation latency. According to Stonebraker and Cetintemel (2005), satisfying the specialized guarantees provided by such solutions might not always be possible, so the workload should be mapped to the appropriate managed service architecture.

Managed services have also started, including modern analytical tools that work on stream models. These analytics can contribute to consistency determinations by potentizing 'hotspots' for user interaction or fault-containing zones. However, these managed services must be integrated to reduce latency or data loss; otherwise, messaging, storage, and analytics pipeline channels must be synchronized.

#### Table 3 Global Consistency Solutions and Their Trade-Offs

Solution	Consistency Model	Key Features	Trade-Offs
Google Spanner	Strong Consistency	<ul> <li>Globally synchronized clock (TrueTime)</li> <li>SQL support with ACID transactions</li> <li>Horizontal scalability</li> </ul>	<ul> <li>Higher latency due to synchronous replication</li> <li>Complex setup and maintenance</li> </ul>
CockroachDB	Strong Consistency	<ul> <li>Distributed SQL with Raft- based replication</li> <li>Automatic failover and scaling</li> <li>ACID transactions</li> </ul>	<ul> <li>Increased latency from consensus protocols</li> <li>Requires careful schema and network design</li> </ul>
DynamoDB	Configurable Consistency	<ul> <li>NoSQL database with flexible consistency settings</li> <li>High availability and low latency</li> <li>Managed service</li> </ul>	<ul> <li>Eventual consistency can lead to temporary data inconsistencies</li> <li>Limited query capabilities compared to SQL</li> </ul>
YugabyteDB	Strong Consistency	<ul> <li>Distributed SQL with high- performance replication</li> <li>Horizontal scalability</li> <li>Multi-region support</li> </ul>	<ul> <li>Complexity in deployment and management</li> <li>Potential latency issues in global deployments</li> </ul>
TiDB	Strong Consistency	<ul> <li>Separation of computation and storage</li> <li>Distributed transactional database</li> <li>Horizontal scalability</li> </ul>	- Requires dynamic scaling configurations - Complex replication protocols
FaunaDB	Strong Consistency	- Globally distributed, serverless database - ACID transactions - Flexible data model	<ul> <li>Potential latency from global replication</li> <li>Vendor-specific features may limit flexibility</li> </ul>
Amazon Aurora Global Database	Strong Consistency	<ul> <li>MySQL and PostgreSQL compatible</li> <li>Global replication with low latency</li> <li>High availability</li> </ul>	<ul> <li>Higher costs for global replication</li> <li>Dependent on AWS infrastructure</li> </ul>

### 14.3. Future Outlook: Balancing Tooling with Architectural Principles

In the future, other frameworks are expected to appear that make the governance of consistency in distributed microservices easier. There are ongoing developments in AI-assisted tools to determine where replicas should be placed and adjust the consistency level in real time based on workload. These innovations present an adaptive solution for overcoming the trade-off between a high degree of consistency and high availability. However, the most important approach still emphasizes the creation of ideal architectures that reflect the business objectives. The seminal work on distributed systems failures and latency by Gray (1978) remains relevant to current advancements in global databases.

Future tooling is expected to come equipped with higher levels of telemetry, where members are learned and replicated factors for latency and the degree of correctness. These capabilities could be especially helpful in tasks with an uncertain workload, while sound architecture, such as well-defined services and an understanding of consistency properties, is essential. Vogels (2009) especially highlighted the need to embrace the idea that it is okay for systems over the scale to converge when scaling systems eventually. However, this decision depended much on the applications' constraints. Distributed system advancements will continue to define how consistency models are addressed. Knowing trade-offs in global consistency solutions, smart use of the managed service, and potential advancements in automation help architects make sound decisions. This approach involves using the early works and adapting the currently trending

systems to each area. Whether in strong consistency, eventual consistency, or a mixture, the vital goal has always been to achieve reliable, high-performing, and customer-oriented solutions in an ever-growing and complicated microservices architecture. All strategies increase the level of hardness and the ability for adaptation.

#### 15. Conclusion

People planning to implement microservices architectures have to weigh the possibilities of using eventual and strong consistency by considering performance, reliability, and scalability. For Eventual consistency, availability and performance are given importance to cater to high throughput and geographical partitioning, but this may lead to some level of data inconsistency that may reduce user confidence in certain applications. On the other hand, strong consistency guarantees that data is up to date and appropriate for use immediately upon being written, which is important in domains such as finance or health care, but this comes at a dear price of high latency and system availability. With the increased integration of distributed systems, these theoretical and practical works have also broadened the options for dealing with these trade-offs. This CAP theorem shows that it is impossible to satisfy the three features: consistency during the normal running of the system. The more recent trends pertain to the further development of consistency models and incorporate AI, adjustability to the microservices requirements and peculiarities, and include such tools as CRDTs and quorum mechanisms.

It is also necessary to highlight the critical part of domain-driven design (DDD) and event-driven architectures in managing consistency models and business goals. In this way, organizations create a more flexible pattern of consistency control based on bounded context and non-synchronous communication patterns of interaction. For example, e-commerce sites combine strong consistency when making payments with eventual consistency choices to update the inventory. Examples from the real world are described in terms of how the consistency choice influences the system development. Transaction protection guarantees are required for financial services to hold their clients to standard to meet data validity and regulatory needs. On the other hand, social media applications focus on availability, employing every consistency level to deal with a high traffic flow to the application. Such examples show that knowledge and context-dependent consistency policies will help systems deliver on operational and user requirements successfully. The paper also considers the issues involved with each approach. Eventual consistency causes users to be confused by inconsistent data and creates dispute resolution challenges for temporary data separation. In contrast, strong consistency relies on synchronous operations, whose performance is worsened by latency and scalability. Patterns such as Saga and Outbox minimize these challenges by creating solid frameworks for event propagation. Some new opportunities for further development include AI adaptive consistency and globally ordered databases such as Google Spanner or Cockroach DB. These tools allow workloads and user requirements to be adjusted for consistency levels in a versatile, fault-tolerant framework for distributed systems.

The consistency model to support should be decided in light of the overall business strategy, technical platform, and users 'expectations. Ongoing refresh is the best practice as the demands and skills required for managing data change over time and with the availability of new tools and frameworks. Organizations can improve their microservices by encouraging training and focusing on architectural rules to meet existing requirements and potential flexibilities. It is important to note that the choice between eventual and strong consistency is not an either-or proposition but rather a question of balance based on the priorities and limitations of the system in question. That is why organizations can develop distributed systems that are reliable, elastic, and coherent with main objectives by recognizing these trade-offs and applying consequent strategies.

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